

Probabilistic Deadline Miss Analysis of Real-Time Systems Using Regenerative Transient Analysis

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RTNS, Versailles - October 10, 2014

Some motivations

- Non-functional requirements prescribed by certification standards e.g., performance requirements on response time and scalability
- Increasing relevance for verification of RAMS requirements: not only Safety, but also Reliability, Availability, Maintainability

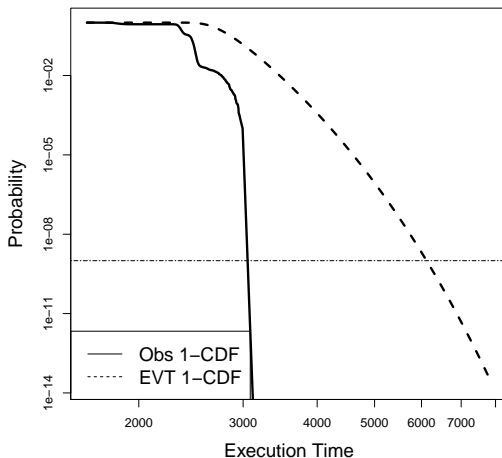
Some areas of investigation

- WCET (Worst Case Execution Time) estimation
 - Static methods (Thesing et al, Healy et al, etc.)
 - Measurement-based methods (Puaut et al, Santinelli et al, etc.)
- Analysis of probabilistic real-time systems
 - Analytical methods (Diaz et al, Cucu-Grosjean et al)
 - State-space based methods (Lindemann et al, Vicario et al, etc.)

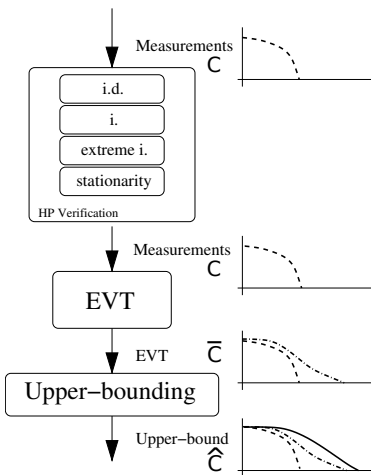
- A probabilistic approach for the analysis of real-time systems with stochastic parameters estimated from real measurements
- Combines results consolidated in the areas of:
 - Estimation of probabilistic WCETs (Extreme Value Theory)
 - Probabilistic analysis (method of stochastic state classes)
- Experimented on a case study including probabilistic WCETs estimated from benchmarks and real system executions
- Targeted to the evaluation of the probability of deadline miss

Probabilistic Worst Case Execution Time (pWCET)

- Given C_k the distribution of the execution time of a task measured in a certain configuration / condition k , the probabilistic Worst-Case Execution Time distribution \bar{C} is an upper-bound on the execution time distribution C_k of all the possible execution conditions of the task



Probabilistic WCET estimation



- Check *identical distribution* (KS test)
 - If OK, verify subsequent hypothesis
- Check *independence* (AR test)
 - If KO, check *stationarity* (LB test) or *independence of extremal samples* (extremogram / extremal index)
- If hypothesis verification is successful, apply the Extreme Value Theory (EVT) to derive a Generalized Extreme Value (GEV) distribution (Gumbel / Fréchet / Weibull)
- Derive the tightest Erlang distribution \hat{C} that upper-bounds the GEV distribution \bar{C}

Probabilistic WCET estimation: experimental results

- Environment for tasks implemented from the Mälardalen benchmark¹
 - A machine with 2 Intel® Xeon® E5620 2.4 GHz sockets each machine having 4 cores and 3 cache levels
 - SchedMCore², for precise real-time execution
 - LTTng (Linux Trace Toolkit new generation)³, for performance monitoring
- Results of hypothesis verification

Task	KS	LB	AR	$\hat{p}(5)$	θ
	(Identical distribution)	(Independence)	(Stationarity)	(Extremogram)	(Extremal index)
<i>ns</i>	0.90	<i>AR</i> (40)	0	0.005	0.995
<i>cnt_isol</i>	0.56	<i>AR</i> (9)	0.0005	0.003	1
<i>cnt_mc</i>	0.26	<i>AR</i> (38)	0	0.005	1
<i>edn</i>	0.82	<i>AR</i> (39)	0	0.01	1

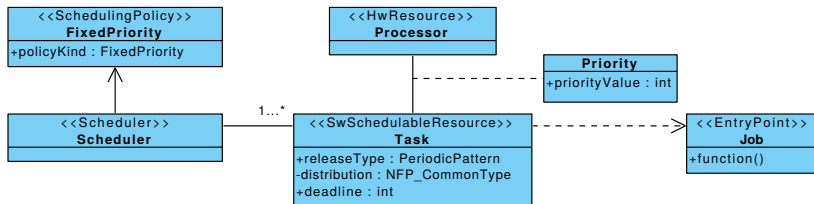
- The algorithm for pWCET estimation can be safely applied

¹ J. Gustafsson, A. Betts, A. Ermedahl, and B. Lisper, The Mälardalen WCET Benchmarks – Past, Present and Future, WCET2010, 2010.

² <https://forge.onera.fr/projects/schedmcore>

³ P. Fournier, M. Desnoyer, and M. R. Dagenais, Combined Tracing of the Kernel and Applications with LTTng, Linux Symposium, 2009.

Task model

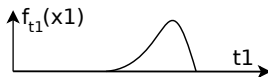
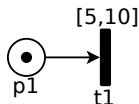


- Single-processor real-time system
- Fixed-priority non-preemptive scheduling policy
 - A higher number corresponds to a higher priority level
- Periodic real-time tasks
 - each associated with a relative deadline
 - each having either a *deterministic* or a *probabilistic* execution time (specified by a *non-Markovian* probability distribution function)
- A job is discarded as soon as its deadline is missed
- Jobs do not use mutex semaphores to synchronize
- A job cannot self-suspend before its completion

Modeling through stochastic Time Petri Nets

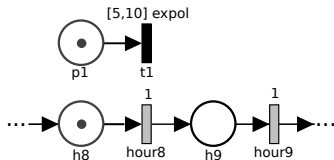
- A class of non-Markovian Stochastic Petri Nets (NMSPN)⁴

- Encompass concurrent GEN transitions with bounded support
- State = marking + remaining times of GEN transitions



- Extended with enabling functions, flush functions, and priorities

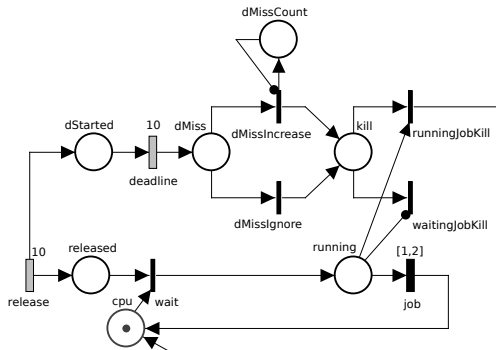
- Change the enabling condition of transitions and the token moves
- Neither restrict the model expressivity nor impact on the analysis



Transition	Enabling function	Flush function	Prio
t1	$h9 == 1$	$\{p4, p5\}$	-
hour8	-	-	10
hour9	-	-	10

Stochastic model of a task

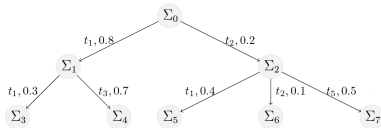
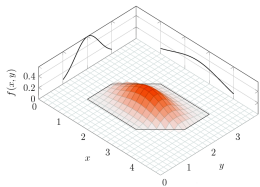
- The underlying process is a Markov Regenerative Process (MRP) including multiple concurrent generally distributed (GEN) timers
 - Regeneration point: the future is independent from the past
 - The model regenerates at each hyper-period
- Example: period=10; deadline=10; execution time supported over [1,2]



Transition	Enabling function	Flush function	Priority
deadline	-	-	9
release	-	-	8
wait	-	-	3
job	-	{dStarted}	-
...

Solution technique: the method of stochastic state classes

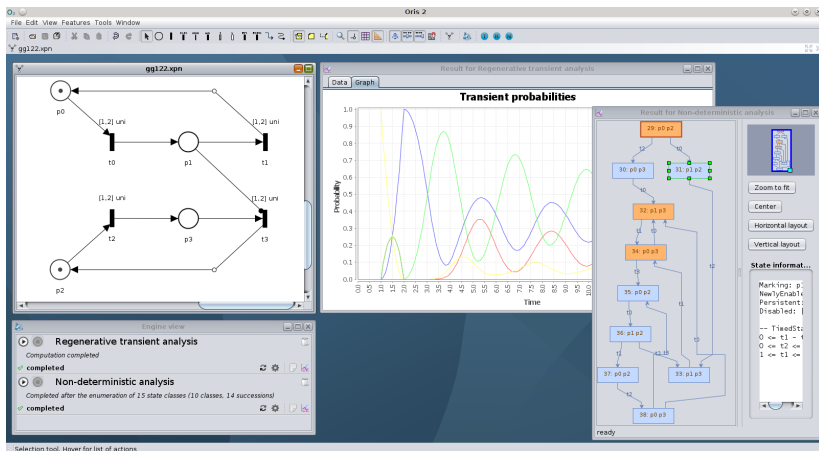
- A *stochastic state class* is a tuple $\langle m, D, f \rangle$
 - $m : P \rightarrow \mathbb{N}$ is a marking
 - $D \subseteq \mathbb{R}_{\geq 0}^n$ is a set of values for times-to-fire $\vec{\tau}$
 - $f : D \rightarrow [0, 1]$ is the PDF of the random vector $\vec{\tau}$ over D
- The successor $\Sigma' = \langle m', D', f' \rangle$ of Σ through transition t holds all the possible states after the firing of t in Σ and their joint PDF
 - D is a Difference Bounds Matrix zone (DBM zone)
 - f is a continuous function (piecewise over DBM subdomains)
- Symbolic calculus implemented in ORIS for expolynomials
- Classes allow to derive kernels of the underlying MRP^{5 6}
 - The transient probability of reachable markings is derived through numerical solution of generalized Markov renewal equations



⁵ V. G. Kulkarni, Modeling and analysis of stochastic systems, CRC Press, 1996

⁶ A. Horváth, M. Paolieri, L. Ridi, E. Vicario. "Transient analysis of non-Markovian models using stochastic state classes". Performance Evaluation, Vol. 69, No. 7, pp. 315â335, July 2012.

The ORIS Tool



- Graphical Petri net editor
- Transient analysis of non-Markovian stochastic Petri nets
- Full Java implementation (cross-platform)
- Available at <http://oris-tool.org>

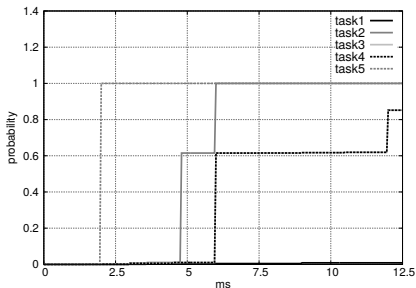
Experimental setting

- A task-set made of 5 periodic tasks
 - 2 tasks have a *deterministic* WCET
 - 3 tasks have a *probabilistic* WCET

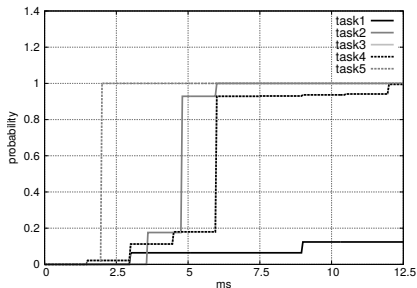
Task	Period	Priority	Execution time	Benchmark
tsk_1	1.0	5	DET(0.2)	-
tsk_2	1.2	4	DET(0.4)	-
tsk_3	1.5	3	Erlang(3, 120.19)	<i>ns</i>
tsk_4	2.0	2	Erlang(3, 40.27)	<i>cnt_isol</i>
tsk_5	3.0	1	Erlang(4, 59.50)	<i>cnt_mc</i>

- Analysis repeated for different distributions of tsk_3 execution time
 - Erlang(2, 50.05), *ns* benchmark
 - Erlang(2, 28.70), *cnt_mc* benchmark
 - Erlang(2, 18.95), *cnt_isol* benchmark
 - Erlang(2, 10.67), *edn* benchmark
- $\sim 15''$ to enumerate state-space, $\sim 30'$ to solve renewal equations

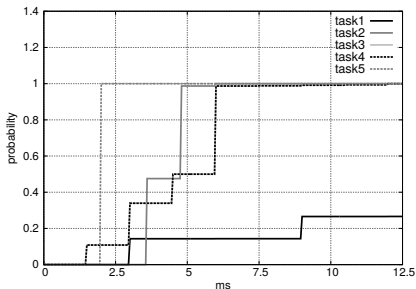
Experimental results: deadline miss probability within time t



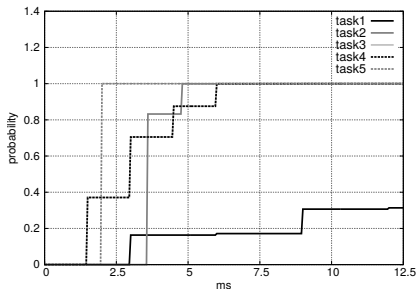
(a) $task_3$ pWCET: Erlang(2, 50.05)



(b) $task_3$ pWCET: Erlang(2, 28.70)



(c) $task_3$ pWCET: Erlang(2, 18.95)



(d) $task_3$ pWCET: Erlang(2, 10.67)

- A probabilistic approach for the analysis of real-time systems
 - Estimation of pWCETs through the Extreme Value Theory
 - Probabilistic analysis through the method of stochastic classes
 - Experiments with pWCETs derived from real benchmarks
- Future issues
 - How do guarantees of pWCETs reflect on schedulability results?
 - What is the accuracy attained by schedulability results if different distributions are used in pWCET estimation?
 - Can other scheduling policies be encompassed in the approach?